

Discrete Mathematics

Logic and Proof

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Logical Equivalence

Tautology and Contradiction

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Exercise. Is $(p \rightarrow q) \leftrightarrow (q \rightarrow p)$ a tautology, a contradiction, or neither?

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- › We write $p \equiv q$ when p and q are logically equivalent.
- › **Example.** $(p \rightarrow q) \equiv (\neg p \vee q)$

Exercise

› Show that

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

Commonly Used Equivalences I

- › Idempotent laws: $p \vee p \equiv p$ $p \wedge p \equiv p$
- › Associative laws:
 $(p \vee q) \vee r \equiv p \vee (q \vee r)$
 $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
- › Commutative laws: $p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$
- › Distributive laws:
 $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$
 $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
- › Identity laws:
 $p \vee F \equiv p$ $p \vee T \equiv T$
 $p \wedge F \equiv F$ $p \wedge T \equiv p$

Commonly Used Equivalences II

- › Double negation law: $\neg\neg p \equiv p$
- › Complement laws: $p \vee \neg p \equiv T$ $\neg T \equiv F$
 $p \wedge \neg p \equiv F$ $\neg F \equiv T$
- › De Morgan's laws: $\neg(p \vee q) \equiv \neg p \wedge \neg q$
 $\neg(p \wedge q) \equiv \neg p \vee \neg q$
- › Absorption laws: $p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$
- › Conditional identities: $p \rightarrow q \equiv \neg p \vee q$
 $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$

Exercise:

› Prove the following De Morgan's law:

$$\neg(p \wedge q) \equiv (\neg p) \vee (\neg q)$$